Next In Line, Please!

Exploiting the Indirect Benefits of Inlining by Accurately Predicting Further Inlining

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Exploiting the Indirect Benefits of Inlining by Accurately Predicting Further Inlining

Efficiently Estimating and Exploiting the Indirect Benefits of Inlining

Master-Thesis von Jannik Jochem
Mai 2011

Fachbereich Informatik
Fachgebiet Softwaretechnik

Background: The Direct and Indirect Benefits of Inlining

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Direct Benefits

- No stack frame creation
- No call/return overhead
- (Possibly) no dynamic dispatch

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- Further, guardless Inlining

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Background: When is Further, Guardless Inlining Possible?

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Precise Arguments

```
class A1 {
   void m() {
      B b = ...
      C c = new D();
      // Precise type of argument is D.
      b.n(c);
   }
}

class D extends C { ... }
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class D extends C { ... }
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Extant Arguments

```
class A2 {
   void m(C c) {
      B b = ...
      // Argument exists before call
      // to A2.m(C).
      b.n(c);
   }
}
```

Background: When is Further, Guardless Inlining Possible? And When Obviously Not?

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Precise Arguments

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   }
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```

```
class D extends C { void n(C c) { this.f = c; } }
```

Background: The Inline Oracle of Jikes RVM

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Problem: Jikes RVM's Assumed Size Reductions

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	Reduction
Reference argument of precise type	15 %
Reference argument pre-exists method call	5 %
Non-null object constant	10 %
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Proposed Solution: Award Size Reductions Only When Further Inlining Likely

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	Reduction
Reference argument of precise type	0 %
Reference argument pre-exists method call	0 %
Non-null object constant	10 %
null constant	10 %
Integer constant	5 %
Array argument of precise type	5 %
No aastore check required	2 %

Proposed Solution: How to Determine Likelihood of Further Inlining?

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Proposed Solution: Identify Precise- and Extant-Induced Edges

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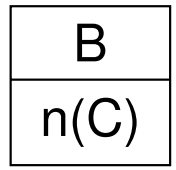


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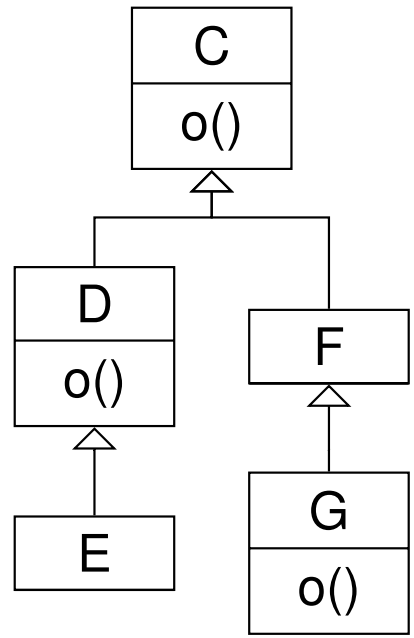

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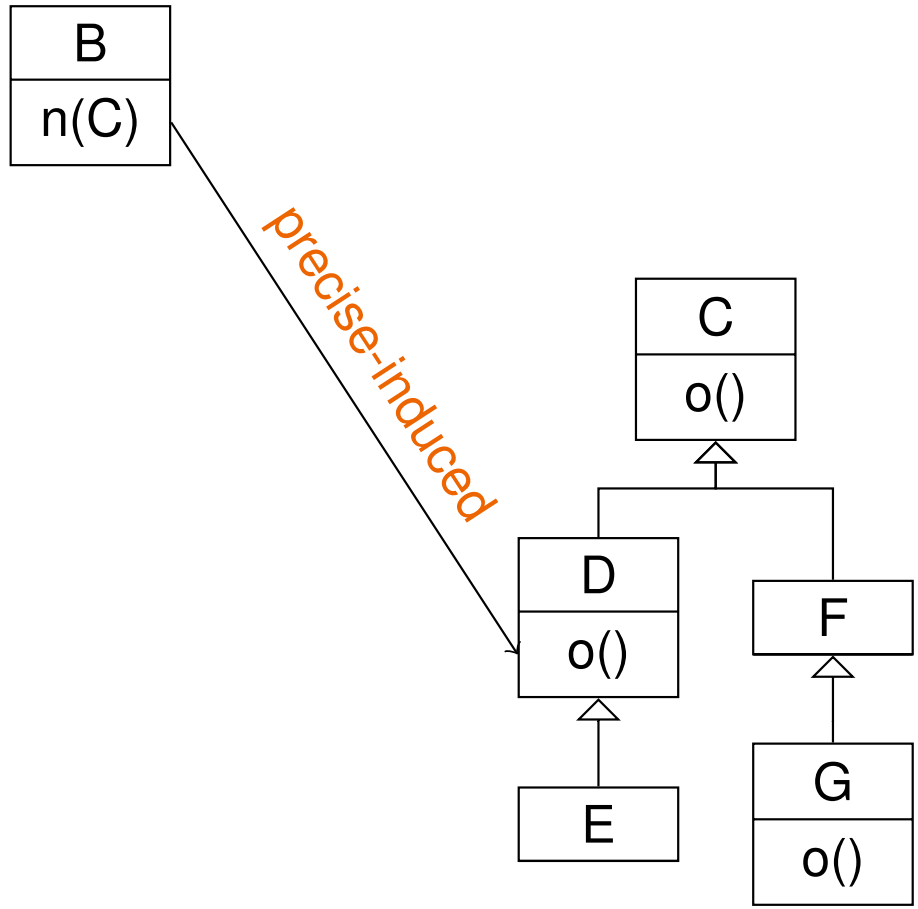

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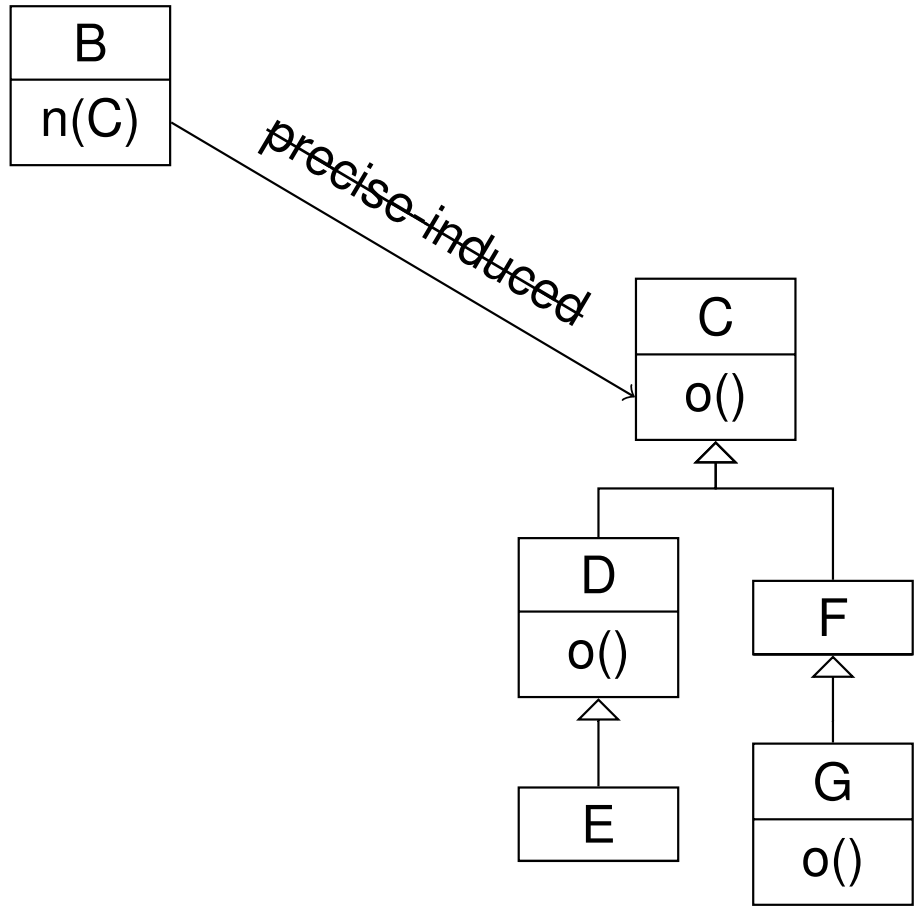

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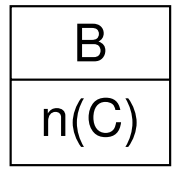


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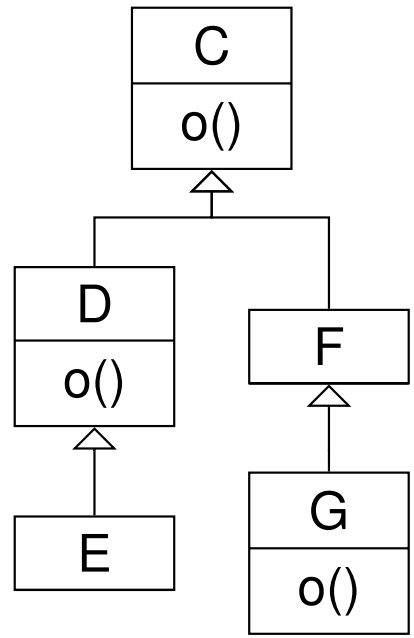
Extant-induced edges work similar.



Evaluation: The Prediction Problem of Further Inlining

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Problem

Given the decision to inline B.n(C) into A.m(), predict whether further inlining of C.o() into B.n(C) will occur—and only then award a size reduction.

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Evaluation: Per-Decision Quality of Inlining Heuristics (Precision)

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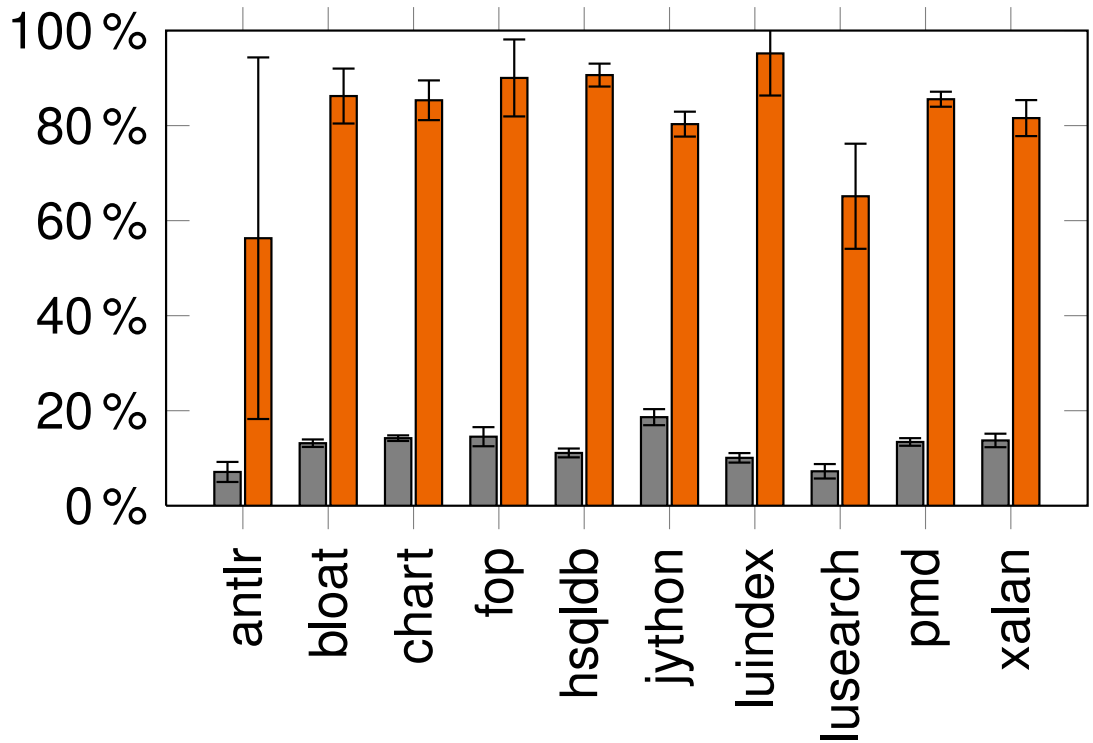

$$\frac{\text{true positives}}{\text{true positives + false positves}}$$

Evaluation: Per-Decision Quality of Inlining Heuristics (Recall)

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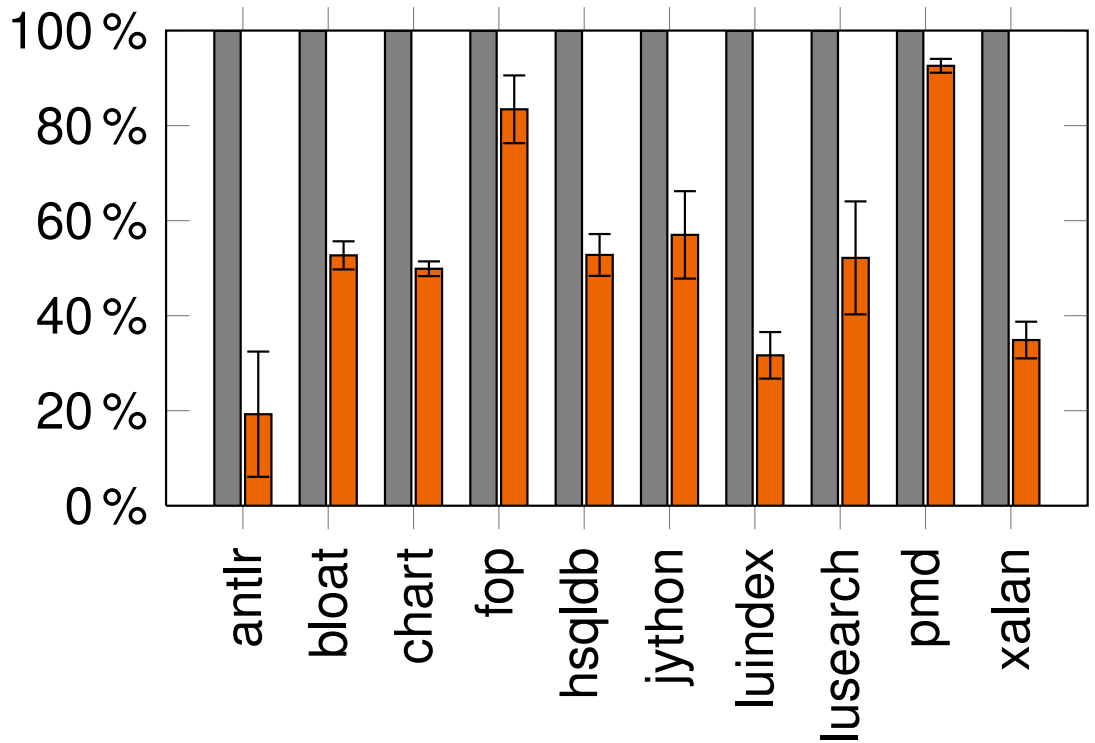

$$\frac{\text{true positives}}{\text{true positives + false negatives}}$$

Evaluation: Per-Decision Quality of Inlining Heuristics (F1-Measure)

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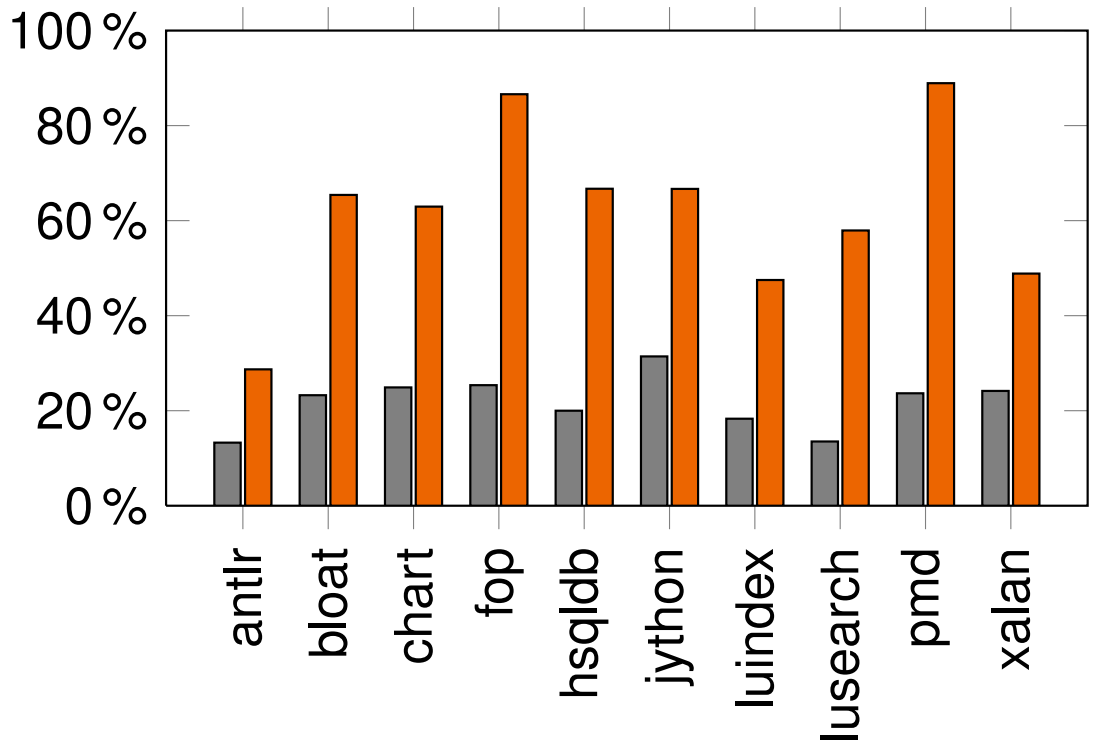

$$2 \times \frac{\text{Precision} \times \text{Recall}}{\text{Precision} + \text{Recall}}$$

Evaluation: Performance Measurements with Replay-Compilation

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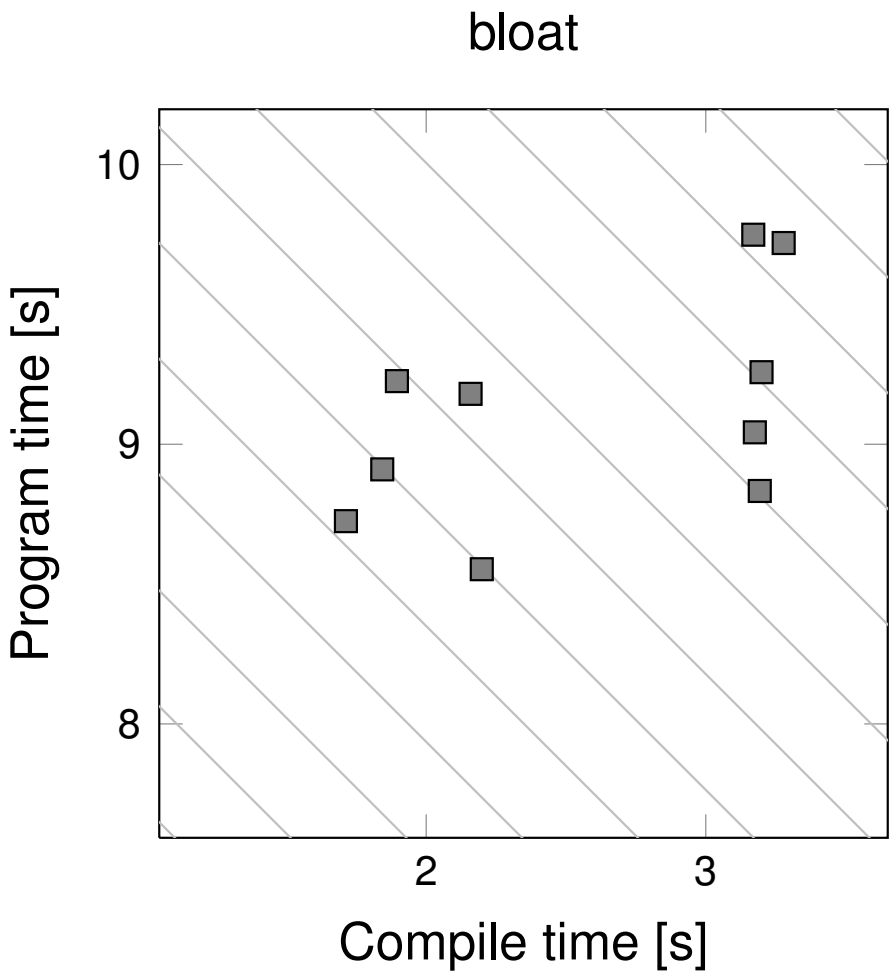

- 10 Compilation Plans

Evaluation: Performance Measurements with Replay-Compilation

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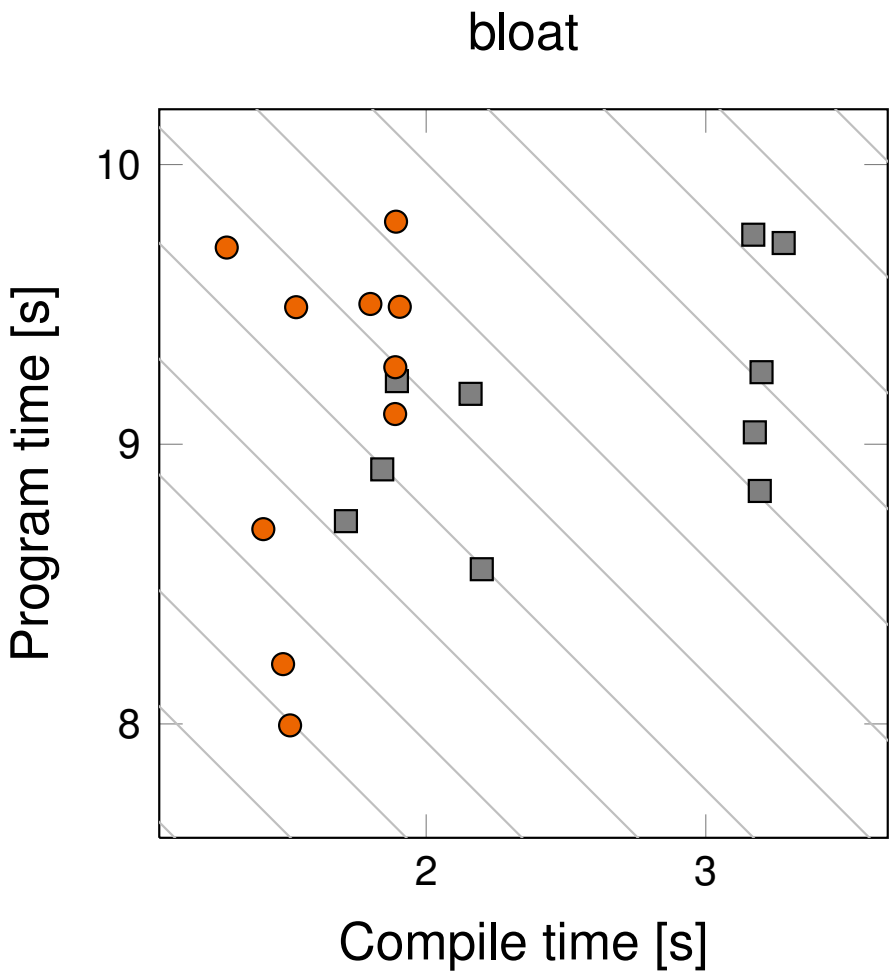

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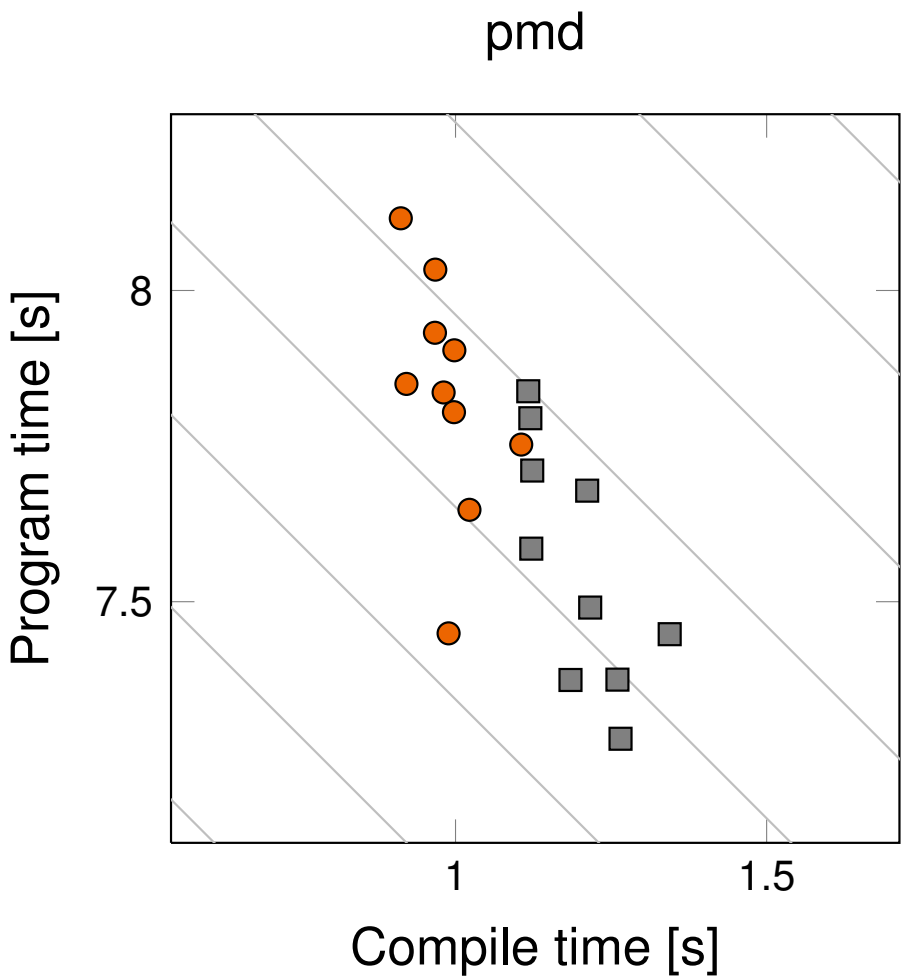

- 10 Compilation Plans
- Proposed Inlining Heuristic saves compile time
- Proposed Inlining Heuristic (sometimes) increases program time

Evaluation: Performance Measurements with and without Replay-Compilation

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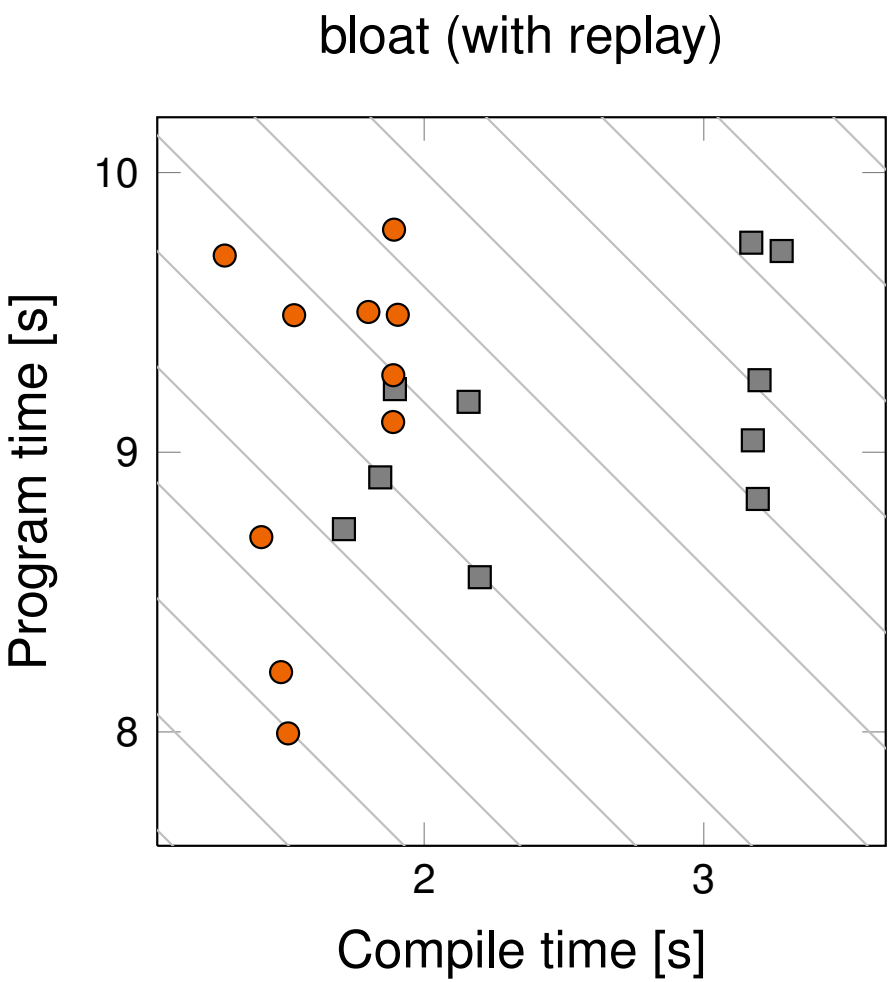

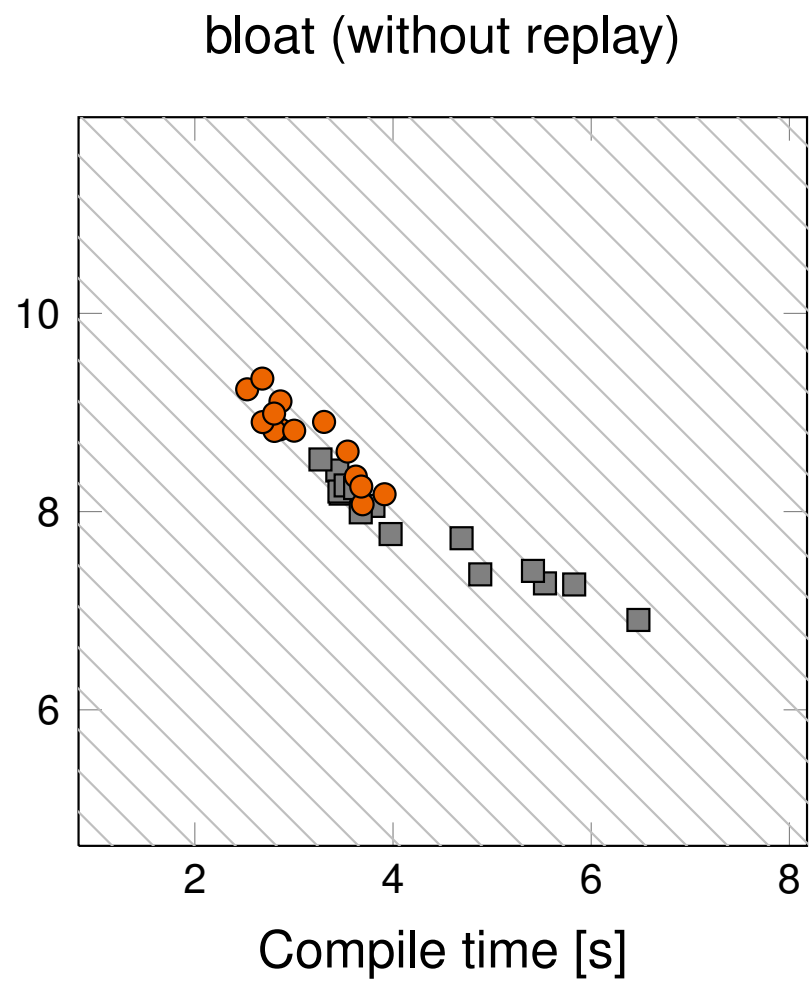

Evaluation: Speed-up on different architectures (AMD Athlon 64, Intel Core i7)

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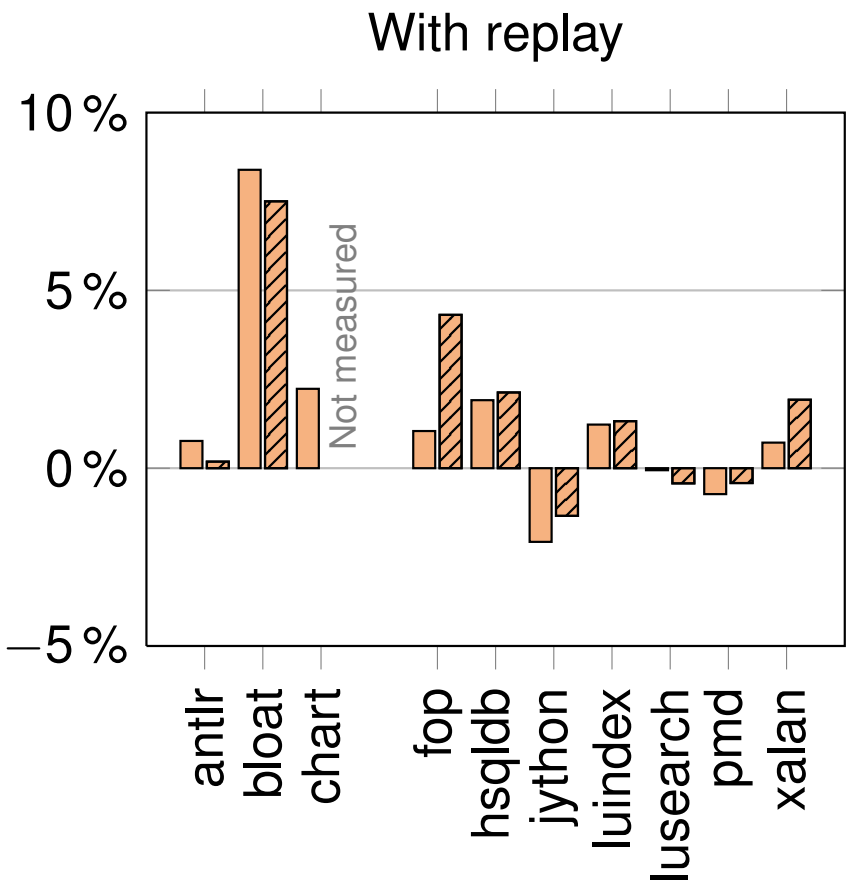

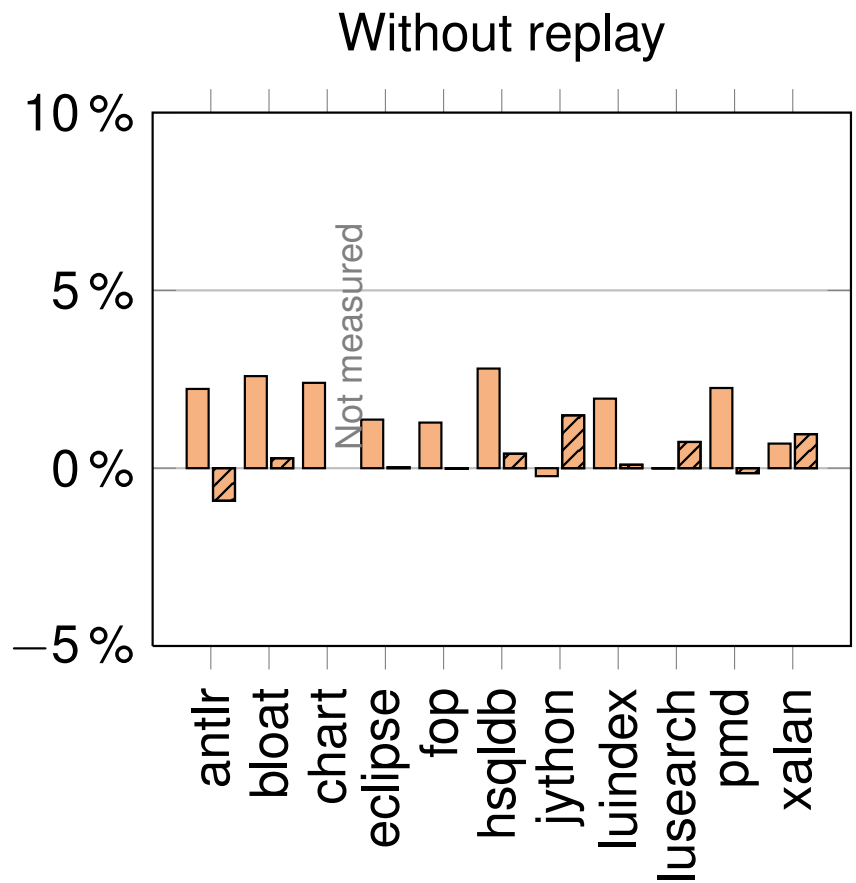

Open Questions and Future Work

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```
Collections.sort(list, new Comparator() {
  int compare(Object lhs, Object rhs) {
    ...
  }
});
```