A Decision Tree-based Approach to Dynamic Pointcut Evaluation

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- Motivation: Dynamic PCD Evaluation
- Approach: Decision-tree based Matching
- Technical Contributions:
 - Formalization of the PCD Evaluation problem
 - Algorithms using Decision-tree structures for faster matching
 - Use of implication relationships for partial evaluation of type predicates

- $a \in \mathcal{A}$, the set of attributes
- $o \in \mathcal{O}$, the set of operators
- $oldsymbol{v} \in oldsymbol{\mathcal{V}},$ the set of values

```
m{a} \in \mathcal{A}, the set of attributes m{o} \in \mathcal{O}, the set of operators m{v} \in \mathcal{V}, the set of values
```

$$\begin{array}{ll} \textit{pred} & ::= & (\textit{\textbf{a}}, \textit{\textbf{o}}, \textit{\textbf{v}}) \\ \textit{fact} & ::= & (\textit{\textbf{a}}, \textit{\textbf{v}}) \end{array}$$

```
egin{array}{lll} {m a} & \in & {\mathcal A}, 	ext{ the set of attributes} \ {m o} & \in & {\mathcal O}, 	ext{ the set of operators} \ {m v} & \in & {\mathcal V}, 	ext{ the set of values} \end{array}
                           pred ::= (a, o, v)

fact ::= (a, v)

        PCD
        ::=
        pred

        |
        ( PCD )

        |
        pred && PCD

        |
        pred || PCD

        join point
        ::=
        fact

        |
        fact && join point
```

```
\mathcal{A} ::= \{ \textit{modifier}, \textit{type}, \textit{name} \}
\mathcal{V} ::= \{ \textit{v} : \textit{v} \text{ is a modifier, type or name in the program} \}
\mathcal{O} ::= \{ ==, != \}
```

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```

Example PCD

```
(modifier, ==, public) && (type, !=, void) &&
(name, ==, "Set")
```

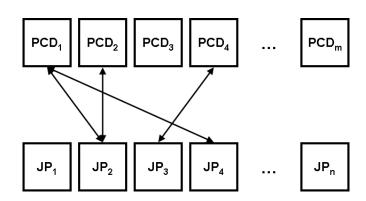
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\mathcal{O} ::= \{ ==, != \}
```

Example PCD

```
(modifier, ==, public) && (type, !=, void) &&
(name, ==, "Set")
```

Example join point

```
(modifier, public) && (type, FElement) &&
(name, "Set")
```



- 2 ways of viewing the problem
 - PCDEval'

PCD1

JP1 JP2 JP3 JP4 ... JPn

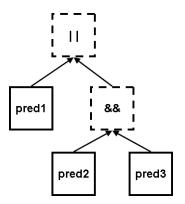
- 2 ways of viewing the problem
 - PCDEval



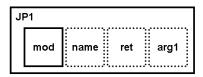


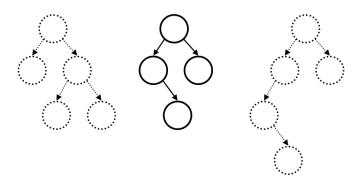
- Evaluation Algorithm overview
 - Order predicates for efficiency
 - Create PCD evaluation tree(s)
 - Add predicates to decision trees
 - Create links to parents

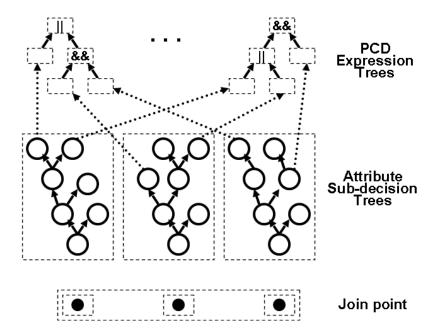
Consider the following PCD: **Pred1**||(**Pred2**&&**Pred3**)

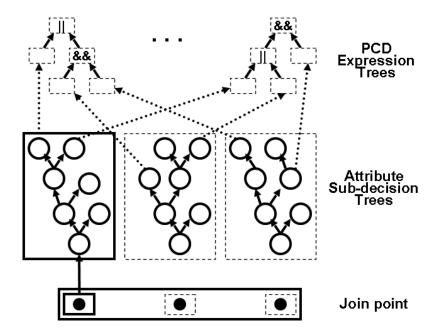


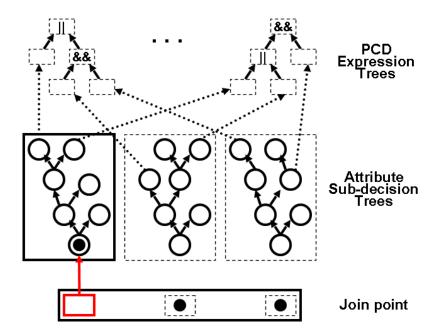
- Order predicates for efficiency
 - Modifiers are simple to match
 - Makes other decision-trees disjoint (smaller)

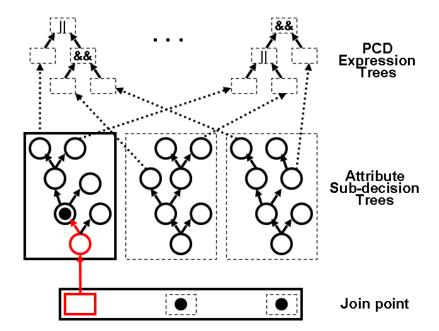


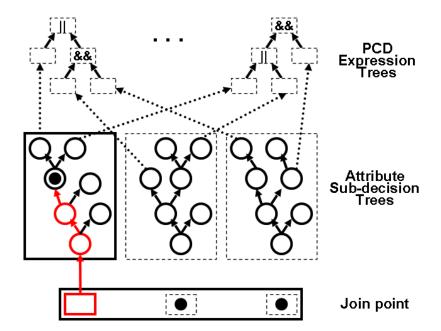


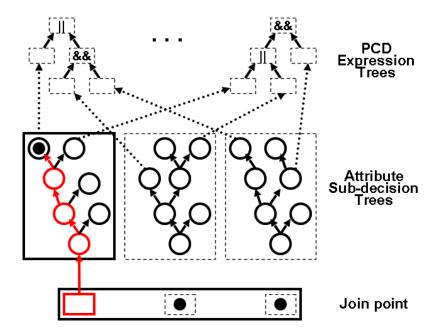


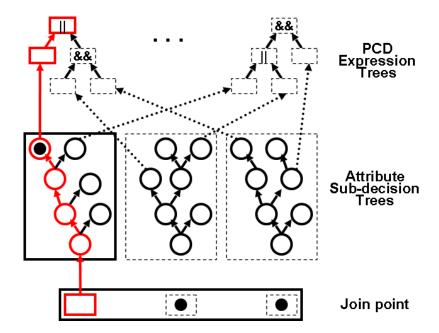


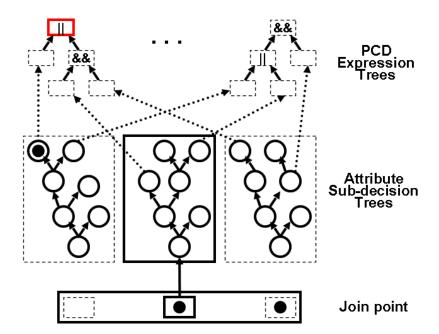


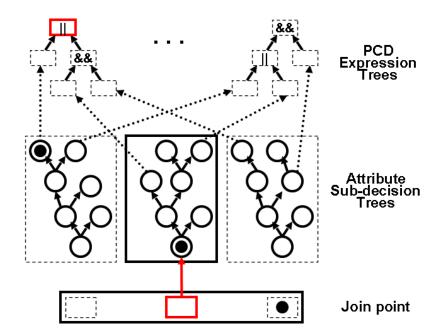


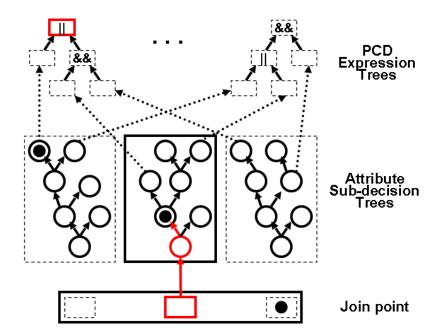


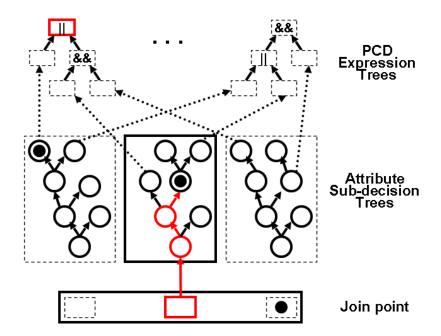


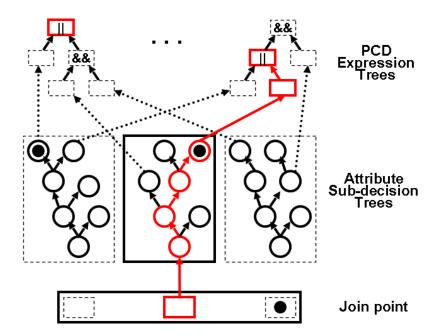


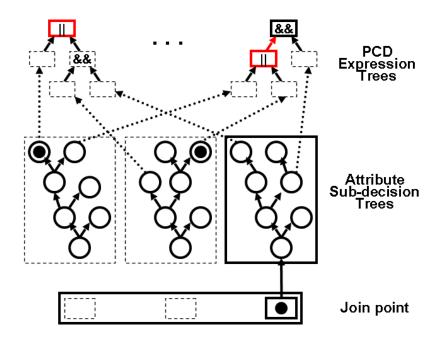


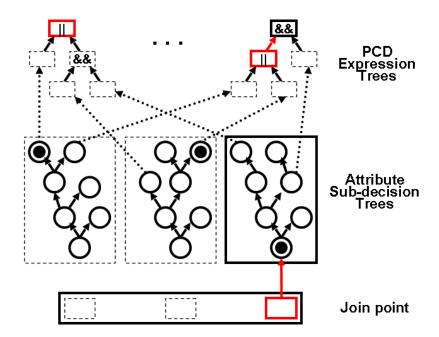


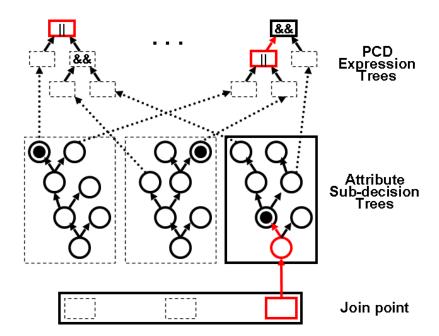


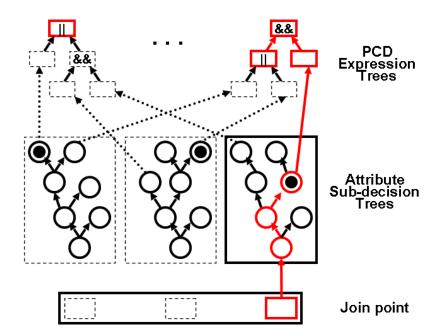


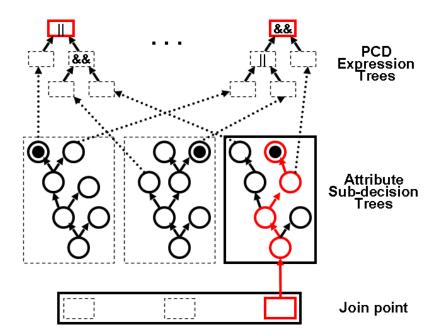












- Goal: Reduce size of decision-trees
- ▶ Idea: Partially evaluate predicates

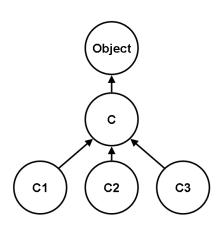
- **▶** Known: *B* < *C*
- Evaluate: **A** ≤ **B**, **A** ≤ **C**

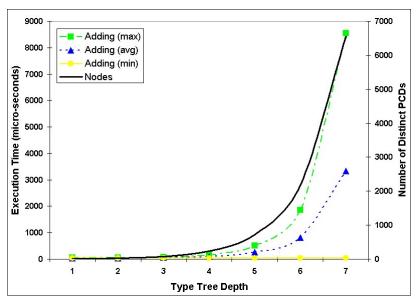
- **▶** Known: *B* < *C*
- Evaluate: **A** < **B**, **A** < **C**
- $ightharpoonup A \lessdot B \land B \lessdot C$

- **▶** Known: *B* < *C*
- Evaluate: **A** ≤ **B**, **A** ≤ **C**
- $\blacktriangleright \ \ \textit{A} \lessdot \textit{B} \land \textit{B} \lessdot \textit{C} \rightarrow \textit{A} \lessdot \textit{C}$

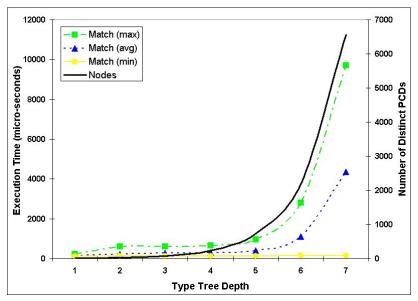
- **▶** Known: **B** < **C**
- Evaluate: **A** < **B**, **A** < **C**
- $\blacktriangleright \ \ \textit{A} \lessdot \textit{B} \land \textit{B} \lessdot \textit{C} \rightarrow \textit{A} \lessdot \textit{C}$
- ▶ Partially Evaluate: A < B</p>

- Created implementation in Nu virtual machine
- Bind and Remove primitives for deploying/un-deploying advice
- Synthetic micro-benchmark
 - Measures time to Bind (add to trees) and match
 - Varies type hierarchy depth





Old matching code - (\sim 40 μ s constant)



Old matching code - average case 3-50x slower worst case 3-88x slower



Related Work

- Efficient Matching Techniques
- Dynamic Residue Evaluation
- Partial Evaluation Techniques

Future Work

- ► Example Implementation(s)
- Real-world Evaluations

- Motivation: Dynamic PCD Evaluation
 - PCDs arrive dynamically
 - PCDs might be removed later
 - Matching the whole (loaded) system against a PCD is too slow
- Approach: Decision-tree based Matching
 - Order evaluations based on cost
 - Partially evaluate wherever possible
- Technical Contributions:
 - Formalization of the PCD Evaluation problem
 - Algorithms using Decision-tree structures for faster matching
 - Use of implication relationships for partial evaluation of type predicates



Overview Problem Our Approach Evaluation Summary

Questions?

http://www.cs.iastate.edu/~nu/

```
C. run
C1. run
C2. run
C3. run
class C {
  public static void run() {
    measure { Bind.. //to methods returning C1 }
    measure { Bind.. //to methods returning C2 }
    measure { Bind.. //to methods returning C3 }
    measure { C1.testMethod }
    measure { C2.testMethod }
    measure { C3.testMethod }
  public C testMethod() { return NULL }
```